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# Detection of Misbehaving Nodes in Ad-Hoc Networks

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Abstract: In wireless ad-hoc networks, a source node relies on intermediate nodes to forward data packets to a designated destination node. Wireless ad-hoc network is deployed in some hostile or uncontrolled environment where nodes may not behave according to the defined protocol. So, there will be a probability of presence of misbehavior nodes. Nodes may participate in route discovery or route maintenance process but refuse to forward packets due to presence of faulty hardware or software or to save their resources, such as, battery power and bandwidth. To avoid such problems, Detection of Node-Misbehavior using Overhearing is introducing. Detection and isolation of misbehavior nodes are important issues to improve the quality of communication service and to save resources of well behaving wireless nodes. In this work, a neighbor Overhearing based Misbehavior Detection (OMD) scheme is proposed. In OMD, each node overhears the transmissions of its neighbors and calculates packet forwarding ratio of its own as well as its neighbors. Source node uses the calculated information to identify a misbehaving node. The proposed scheme reduces communication overheads and identification delays to detect misbehaving nodes in wireless ad-hoc network. In this project, simulation is done by using ns3 software. Simulation results are presented to evaluate the performance of the proposed OMD scheme.

Key words: OMD, Ad-Hoc Networks.

# **1. INTRODUCTION**

Wireless ad-hoc network is a decentralized network which is not supported by any pre-existing infrastructure. Nodes in this type of network work collaboratively to realize end to end communication.

Wireless based Networks have changed today's life greatly, while ad-hoc networks provide people more solutions and convenient due to its special property. Ad hoc networks are mainly formed by a group of devices. The devices are called "nodes" in ad hoc network.

### 2. WORMHOLE ATTACKS:

Wormhole attack is also known as tunneling attack. A As the nodes in MANETs are battery powered, energy tunneling attack is where two or more nodes may becomes a precious resource, and thus, role of selfish collaborate to encapsulate and exchange messages nodes draws more attention. Thus, it is introduced between them along existing data routes. This exploit altogether three routing behaviors of nodes in a MANET. gives the opportunity to a node or nodes to short-circuit Type-0: well-behaved node: A well behaved node the normal flow of messages creating a virtual vertex cut in the network that is controlled by the two colluding attackers. Lack of Co-operation Mobile ad-hoc networks rely on the co-operation of all the participating nodes. The more nodes co-operate to transfer traffic, the more powerful a MANET gets. A selfishness node want to preserve own resources while using the services of others and consuming their resources.

Major contributions of this work are as follows:

• An improved neighbor overhearing based misbehavior detection approach has been proposed to identify misbehaving nodes in the network.

the proposed approaches.

### 3. IMPROVED NODE MISBEHAVIOR DETECTION

### **NODE MISBEHAVIORS:**

Identification of misbehaving nodes in ad hoc networks is critically important to detect security attack in the network. Two types of misbehaving nodes such as selfish and malicious nodes are taken into consideration in [1,2]. Selfish nodes do not intend to directly damage other nodes, but however, do not cooperate, saving battery life for their own communications. But malicious nodes do not give priority to saving battery life, and aim at damaging other nodes. It is introduced that two different types of selfish nodes.

cooperates in the communication well, performs as required by the routing protocol, and equally participates in the communication activities like route discovery, maintenance, packet forwarding and receiving etc.

Type-1: active selfish node: Such a node does not participate in packet forwarding, and drops every received packet. It disables the packet forwarding mechanism for the packets which have a destination address, other than this selfish node. In fact, it helps the selfish node to save its own energy, thereby still contributing to network maintenance.

• Simulation results are presented to show the efficacy of **Type 2**: passive selfish node: Such a node practically does nothing and stays idle in the network. It does not



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contribute to any of the activities like packet forwarding, T1.Destination node calculates PFR of itself and its receiving, route discovery, network maintenance. With neighbors and appends this value to the control packet Req respect to above mentioned misbehaving nodes, we PFR and sends the packet in reverse direction to the source evaluate the performance of DSDV, DSR and AODV node. All intermediate nodes append their calculated PFR routing protocols through extensive simulations, where a value to the packet and send it to their previous node. certain percentage of nodes behave as active and/or Source node updates its matrix based on the values in the passive selfish nodes with the remaining nodes being wellbehaved.

### **OVERHEARING BASED MISBEHAVIOR DETECTION:**

In this work, the proposed neighbor Overhearing Based 4.2 Algorithm: Misbehavior Detection (OMD) technique uses the The following algorithm is used to detect the misbehavior information provided by the neighbors for behavior node: evaluation of nodes in the path without incurring high **Require:** Forwarding ratio PFR of each node, Time epoch communication overhead. An algorithm for overhearing t and t1 based misbehavior detection is shown in (Algorithm guk1) **Ensure:** Detection of misbehaving node Source node maintains a matrix *opp*R of Packet 1. Initialize matrix *opp*R of Packet Forwarding Ratio Forwarding Ratio (PFR) of each node. The matrix contains (PFR) with zeros. node ID denoted as i, forwarding ratio calculated by its 2. Selection of path PS->D using DSR algorithm between own denoted as PFRii and forwarding ratio calculated by source node S and destination node D. its neighbors denoted as PFRij, where j denotes all the 3. Monitoring packet delivery ratio n of the path PS->D neighbors of node ni.

Each node calculates PFR of its own and its neighbors and 4. If  $\eta < \eta 0$ these values are used to update the matrix  $\varphi$ PFR of source 5. Node S sends a control packet ReqPFR to request all the (refer Algorithm).

IPV4 ROUTTING PROTOCOL: The base class defines and its neighbors PFR<sub>ii</sub> (where j denotes all the neighbors two packet routing and forwarding. AODV discovers routes on demand. Therefore, the AODV model buffers all 6. Calculation of PFRii and PFRji by the intermediate packets while a route request packet is disseminated. The node. AODV implementation can detect the presence of 7. At time t+t1 node D sends a reply RepPFR along the unidirectional links and avoid them if necessary. When reverse path and all intermediate nodes append the source node wants to send a packet to some destination calculated values to the reply. node, it uses the Dynamic Source Routing (DSR) 8. Node S modifies it's matrix  $\varphi$ PFR based on the received algorithm to find the route to the destination node. In RepPFR DSR, when source node has a packet for destination node, 9. Node S checks all the entries it checks whether a route exists in its cache.

If route exists in the route cache, source node uses the 11. Node i is misbehaving node. route to send packets to the destination node. If a route 12. Else does not exist, source node broadcasts a Route Request 13. Node i is wellbehaving node. (RREQ) packet. This packet contains the source ID, destination ID, and the Time-To-Live (TTL). Any 4. SIMULATION RESULTS intermediate node that receives the RREQ, appends its ID to the RREQ packet and rebroadcasts it while decreasing **4.1 CREATION OF NODES** the TTL field by one unit. If a receiving node is the destination node, it responds to source node with a Route Reply (RREP) packet containing the entire path PS->D from source node to destination node.

Source adds the route to route cache for further use. Source node sends the data packets to the destination node using the route determined. Source node continuously monitors the performance of path PS->D based on the packet delivery ratio n. If it detects that this ratio is less than some predefined threshold value  $\eta 0$  (refer Algorithm), it will send a control packet Req PFR at time 't' to request all the intermediate nodes to calculate Packet Forwarding Ratio (PFR) of their own and their neighbors' based on the number of packets forwarded Ps and number of packets received Pr for a time duration T1.

When this control packet reaches to the destination node, it generates a control packet RepPFR after time duration

control packet Rep PFR. Source node checks the matrix and finds the node for which calculated values of Packet Forwarding Ratio conflict with each other and declares that node as misbehaving node.

by node S.

nodes along the path to calculate PFR of its own PFRIii of node i) at time t for a time duration t1.

10. If PFRii\_= PFRij.



Figure 1: Creation of nodes (red dot indicates nodes).



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4.2 CASE 1: HIDDEN PASSIVE ATTACK (100% 4.3 PACKET LOSS) (100

4.3 CASE2: EXPOSED PASSIVE ATTACK (100%PACKET LOSS)







Simulating Case:(1) Creating 22 nodes Encountering 100% Packet Loss Due To Hidden Wo Starting simulation for 35 s ...

Figure 4: 100% packet loss in Hidden passive attack .



Figure 5: Misbehaving node is exposed with 100% packet loss.

<pre>gobject.sonre_remove(setrupdate_timeout_to) ser@user.VirtualBox:~/ns-allinone-3.22/ns-3.22\$ ./wafrun "scratch/project1cas laf: Entering directory `/home/user/ns-allinone-3.22/ns-3.22/build' laf: Leaving directory `/home/user/ns-allinone-3.22/ns-3.22/build' build' finished successfully (1.985s) imulating Case:(2) ireating 22 nodes</pre>
laf: Entering directory `/home/user/ns-allinone-3.22/ns-3.22/build' laf: Leaving directory `/home/user/ns-allinone-3.22/ns-3.22/build' build' finished successfully (1.985s) imulating Case:(2) ireating 22 nodes
af: Leaving directory `/home/user/ns-allinone-3.22/ns-3.22/build' build' finished successfully (1.985s) imulating Case:(2) ireating 22 nodes
build' finished successfully (1.985s) imulating Case:(2) reating 22 nodes
imulating Case:(2) ireating 22 nodes
reating 22 nodes
ctive Powerful Wormhole
tarting simulation for 35 s
5:21:19 environ No en_IN translation found for domain kiwi
ould not load icon applets-screenshooter due to missing gnomedesktop Python module
canning topology: 22 nodes
canning topology: calling graphviz layout
canning topology: all done.
ING 10.0.0.22 56(84) bytes of data.
home/user/ns-allinone-3.22/ns-3.22/src/visualizer/visualizer/core.py:1098: Warning:
found when attempting to remove it
gobject.source_remove(self.update_timeout_id)
10.0.0.22 ping statistics
5 packets transmitted, 0 received, 100% packet loss, time 34999ms
home/user/ns-allinone-3.22/ns-3.22/src/visualizer/visualizer/core.py:1113: Warning:
found when attempting to remove it
gobject.source_remove(self. update_timeout_id)
<pre>goojeec.source_remove(ser:_update_traoue_co) ser@user.VirtualBox:~/ns-allinone-3.22/ns-3.22\$</pre>

Figure 6: Simulation result for case 2.



Figure 7: Output terminal shows 100% packet loss in exposed passive attack.



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Vol. 3, Issue 6, June 2016 4.4 CASE 3: HIDDEN ACTIVE ATTACK (CERTAIN

% PACKET LOSS)



200

00		-100			0	100	200
Zoom:	0.517	\$ Speed:	1.000	\$	Time: 1.200000 s	Snapshot	Simulate (F3)
Advar	nced						



•200 00 -100 0 100 200 Zoom: 0.517 ↓ Speed: [5.300 ↓ Time: 90.650000 s Snapshot Simulate (F3)

Advanced

Figure 8: Reduction of packet loss by 18 percent.

user@user-VirtualBox:~/ns-allinone-3.22/ns-3.22\$ . Waf: Entering directory `/home/user/ns-allinone-3.2 laf: Leaving directory `/home/user/ns-allinone-3.22 'build' finished successfully (2.059s) Simulating Case:(3) Creating 22 nodes Starting simulation for 35 s ... 15:23:39 environ No en IN translation Could not load icon applets-screenshooter due to mis scanning topology: 22 nodes... scanning topology: calling graphviz layout scanning topology: all done. PING 10.0.0.22 56(84) bytes of data. 64 bytes from 10.0.0.22: icmp\_seq=0 ttl=62 time=57 /home/user/ns-allinone-3.22/ns-3.22/src/visualizer/v found when attempting to remove it gobject.source\_remove(self.\_update\_timeout\_id) /home/user/ns-allinone-3.22/ns-3.22/src/visualizer/v found when attempting to remove it gobject.source\_remove(self.\_update\_timeout\_id) --- 10.0.0.22 ping statistics ---35 packets transmitted, 6 received, 82% packet loss, /home/user/ns-allinone-3.22/ns-3.22/src/visualizer/v t found when attempting to remove it gobject.source\_remove(self.\_update\_timeout\_id) /home/user/ns-allinone-3.22/ns-3.22/src/visualizer/v found when attempting to remove it gobject.source\_remove(self.\_update\_timeout\_id) user@user-VirtualBox:~/ns-allinone-3.22/ns-3.22\$ Figure 9: Simulation result for case 3.



0: Output terminal shows reduced packet los 18 percent.



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# 4.5 CASE 4: HANDLES WITH RATE ADAPTION the network. From the experience of this, the following ALGORITHM (LESS PACKET LOSS)



### Simulating Case:(5)

Creating 22 nodes Starting simulation for 35 s ... 16:09:44 environ No en IN translation found for domain ki Could not load icon applets-screenshooter due to missing gnomedesktop scanning topology: 22 nodes... scanning topology: calling graphviz layout scanning topology: all done. PING 10.0.0.22 56(84) bytes of data. 64 bytes from 10.0.0.22: icmp\_seq=0 ttl=62 time=55 ms 54 bytes from 10.0.0.22: icmp\_seq=1 ttl=62 time=1 ms j4 bytes from 10.0.0.22: icmp\_seq=2 ttl=62 time=1 ms i4 bytes from 10.0.0.22: icmp\_seq=3 ttl=62 time=1 ms j4 bytes from 10.0.0.22: icmp\_seq=4 ttl=62 time=1 ms 64 bytes from 10.0.0.22: icmp\_seq=5 ttl=62 time=1 ms 64 bytes from 10.0.0.22: icmp\_seq=6 ttl=62 time=1 ms 64 bytes from 10.0.0.22: icmp\_seq=7 ttl=62 time=1 ms 64 bytes from 10.0.0.22: icmp\_seq=8 ttl=62 time=2 ms 64 bytes from 10.0.0.22: icmp\_seq=9 ttl=62 time=2 ms 64 bytes from 10.0.0.22: icmp\_seq=15 ttl=61 time=2008 ms 64 bytes from 10.0.0.22: icmp\_seq=26 ttl=61 time=22 ms 64 bytes from 10.0.0.22: icmp\_seq=31 ttl=62 time=3 ms --- 10.0.0.22 ping statistics ---35 packets transmitted, 17 received, 51% packet loss, time 34999ms rtt min/avg/max/mdev = 1/124.1/2008/485.7 ms

Figure 11: Simulation result for case 4 shows 51% packet loss.

### **5. CONCLUSION**

The frame work for 22 nodes (assumption) was created in ad-hoc networks using NS3 software. Each node is simulated and a program is written for specific nodes to act as malicious node. Using OMD technique, misbehavior nodes were detected and packet forward ratio is determined using simulation table. The OMD technique reduces the identification delay by transmitting few control packets by identifying the misbehaving nodes in

recommendation on future work is suggested: This work can be extended by giving mobility to each node in the network.

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